

# 数理逻辑的程序可靠性验证

顾名宇

561000

摘要:

SPIN

关键词:

中图分类号: TP301

文献标识码: A

文章编号: 1000-2324(2015)04-0621-04

## Validation of Reliability on Mathematical Logic Program

GU Ming-yu

561000,

**Abstract:** The reliability verification of program often takes a very long time to develop the software, while the current software reliability verification method is mainly based on formal methods such as SPIN based on model checking method, but this method may lead to the complexity of verification to improve greatly due to model problems, and then finally fails in validation results everywhere. In order to solve this problem, this paper used mathematical methods from mathematical logic perspective to realize the program reliability verification, and completed the reliability verification of client server program, the method was proved the validity in true.

**Keywords:** Mathematical logic; program; validation

SPIN

[2][6]

100%

90%

70%

## 1 数理逻辑

[5]

1847

### 1.1 命题演算

收稿日期: 2013-07-13

修回日期: 2013-07-24

作者简介: (1960-), , .

E-mail:gumingyu60@126.com

: 2015-03-16 <http://www.cnki.net>

x==0,

y=5;

P: "x==0"

Q: "y=5"

" x==0, y=5" P->Q

### 1.2 谓词演算

x==0, y=5 P(x,y) x==y,Q(x,y) x=y,x y  
P(x,0)->Q(y,5)

## 2 使用数理逻辑验证程序可靠性方法

### 2.1 程序可靠性定义

IEEE

" "

- 1
- 2

### 2.2 程序可靠性验证步骤

1 [5] [3]  
V(xi)(0<=i<=n) xi F1,F2,...Fm x1,x2,x3,...,xn n  
P1,P2,...PK S  
y1,y2,...yl W(yj) yj , 0<=j<=l,  
2  
v v v  
^ ^  
S  
0<=l<=n,0<=t<=n

T " "

3

[1]

S

### 3 用数理逻辑法验证客户服务器程序

R1..Rn

R1.....Rn

R\_Work[Max] Max  
= 0 (0<= i < Max)

R\_Work[i].bz

```

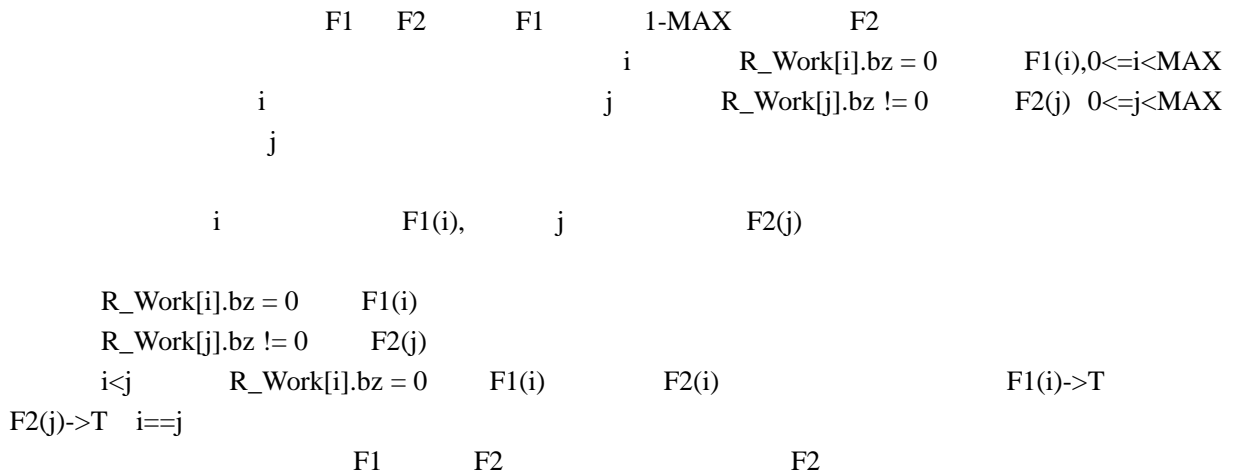
R_Work[Max] F1
for (int i = 0; i < Max; i++)
{
if(R_Work[i].bz == 0)
{
break
}
}

```

```

F2
for int j = 0 j < Max; j++
{
if(R_Work[j].bz != 0)
{
R_Work[j].bz = 0
break
}
}

```



```

F1 F2
R_Work
Max*log Max
R_Nextwork ( 0) F1 F2
程序 F3:
int k;
for (int i = 0; i < Max; i++)
{
k = (R_Nextwork+i) % Max;
if (R_Work [k] .bz == 0)
{
提交任务
break
}
}
程序 F4:
int k
for (int i = 0; i < Max; i++)
{
k = (R_Nextwork+j) % Max;
if (R_Work[k].bz != 0)
{
执行任务
R_Work[k].bz = 0;
R_Nextwork = (k + 1) % Max;
break;
}
}

```

[4] F1 F2

$R_1, \dots, R_n, \quad R_i \ R_k \ R_i < R_k$   
 $F_3(R_i) \rightarrow T \quad F_3(R_k) \rightarrow T \quad F_4(R_k) \rightarrow T$   
 $F_4(R_i) \quad F_4(R_k) \quad , \quad F_4(R_i) \rightarrow T$   
1      F3      F3(R<sub>i</sub>)→T   F3(R<sub>k</sub>)→T    R<sub>i</sub><R<sub>k</sub>      R<sub>i</sub>      R<sub>k</sub>  
2      F4      F4(R<sub>k</sub>)→T      F4(R<sub>k</sub>)      R<sub>i</sub>      R<sub>Nextpoll</sub>  
= R<sub>i</sub>      k    k = (R<sub>Nextpoll</sub> + i)%Max    i > 0    R<sub>k</sub>    R<sub>i</sub>      R<sub>k</sub>      k  
1    2      R<sub>i</sub>    R<sub>k</sub>      ,    F4(R<sub>k</sub>)→T      F4(R<sub>i</sub>)→T  
F4(R<sub>k</sub>)

### 4 小结

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